

Γ	S	H
S	$R + \delta \Gamma$	$\Delta \Gamma$
H	$r + \Delta \Gamma$	$r + \Delta \Gamma$

(a) Analyze the Iterated Stag Hunt with $0 < r < R$ and $0 < \delta < \Delta < 1$.

$S \times S$	$S \times H$	$H \times S$ (or $H \times H$)
$E = R + \delta E$	$E = \Delta E$	$E = r + \Delta E$
$(1 - \delta)E = R$	$(1 - \Delta)E = 0$	$(1 - \Delta)E = r$
$E = \frac{R}{1 - \delta}$	$E = 0$	$E = \frac{r}{1 - \Delta}$

So, the overall payoff matrix (for the symmetric game with payoffs to the row player) is:

Γ	S	H
S	$\frac{R}{1 - \delta}$	0
H	$\frac{r}{1 - \Delta}$	$\frac{r}{1 - \Delta}$

- S is an ESS if $\frac{R}{1 - \delta} > \frac{r}{1 - \Delta}$.
- H is an ESS if $\frac{r}{1 - \Delta} > 0$, i.e. always.

(b) Analyze $\Gamma = (\Gamma_1, \Gamma_2)$ with $0 < r < R$ and $0 < \delta < \Delta < 1$.

Γ_1	S	H
S	$R + \delta \Gamma_2$	$\Delta \Gamma_1$
H	$r + \Delta \Gamma_1$	$r + \Delta \Gamma_1$

Γ_2	rest
rest	$\Delta \Gamma_1$

• The payoffs for the row player in Γ_1 are the same as in part (a) for $S \times H$, $H \times S$ and $H \times H$

• For $S \times S$, $\begin{cases} E_1 = R + \delta E_2 \\ E_2 = \Delta E_1 \end{cases} \Rightarrow \begin{cases} E_1 = R + \delta \Delta E_1 \\ (1 - \delta \Delta) E_1 = R \\ E_1 = \frac{R}{1 - \delta \Delta} \end{cases}$

So, the overall payoff matrix is

Γ	(S, rest)	(H, rest)
(S, rest)	$\frac{R}{1 - \delta \Delta}$	0
(H, rest)	$\frac{r}{1 - \Delta}$	$\frac{r}{1 - \Delta}$

- (S, rest) is an ESS if $\frac{R}{1 - \delta \Delta} > \frac{r}{1 - \Delta}$
- (H, rest) is an ESS if $\frac{r}{1 - \Delta} > 0$, i.e. always.

(c) Analyze $\Gamma = (\Gamma_1, \Gamma_2, \Gamma_3)$ with $0 < r < R$ and $0 < \delta < \Delta < 1$.

Γ_1	S	SR	H
S	$R + \delta \Gamma_1$	$R + \delta \Gamma_2$	$\Delta \Gamma_1$
SR	$R + \delta \Gamma_2$	$R + \delta \Gamma_3$	$\Delta \Gamma_1$
H	$r + \Delta \Gamma_1$	$r + \Delta \Gamma_1$	$r + \Delta \Gamma_1$

Γ_2	rest (SR player)
hunt a hare (S player)	$r + \Delta \Gamma_1, \Delta \Gamma_1$

Γ_3	rest (SR player)
rest (SR player)	$\Delta \Gamma_1$

In the overall payoff matrix, SxS, SxH, HxS and HxH yield the same payoffs as in part (a). SRxH yields the same payoff as SxH. HxSR yields the same payoff as HxS + HxH.

SxSR $(S, \text{hare}) \times (SR, \text{rest})$

$E_1 = R + \delta E_2$ $E_2 = r + \Delta E_1$

$E_1 = R + \delta(r + \Delta E_1)$

$E_1 = R + \delta r + \delta \Delta E_1$

$(1 - \delta \Delta) E_1 = R + \delta r$

$E_1 = \frac{R + \delta r}{1 - \delta \Delta}$

SRxS $(SR, \text{rest}) \times (S, \text{hare})$

$E_1 = R + \delta E_2$

$E_1 = R + \delta \Delta E_1$

$(1 - \delta \Delta) E_1 = R$

$E_1 = \frac{R}{1 - \delta \Delta}$

$(SR, \text{rest}) \times (S, \text{hare})$

$E_2 = \Delta E_1$

yields the same payoff as HxS + HxH.

SRxSR $(SR, \text{rest}) \times (SR, \text{rest})$

$E_1 = R + \delta E_3$ $E_3 = \Delta E_1$

$E_1 = R + \delta \Delta E_1$

$(1 - \delta \Delta) E_1 = R$

$E_1 = \frac{R}{1 - \delta \Delta}$

So, the overall payoff matrix is

Γ	$(S, \text{hare}, \text{rest})$	$(SR, \text{hare}, \text{rest})$	$(H, \text{hare}, \text{rest})$
$(S, \text{hare}, \text{rest})$	$\frac{R}{1 - \delta}$	$\frac{R + \delta r}{1 - \delta \Delta}$	0
$(SR, \text{hare}, \text{rest})$	$\frac{R}{1 - \delta \Delta}$	$\frac{R}{1 - \delta \Delta}$	0
$(H, \text{hare}, \text{rest})$	$\frac{r}{1 - \Delta}$	$\frac{r}{1 - \Delta}$	$\frac{r}{1 - \Delta}$

where the conditional strategy (a) means "Play a if the row player in Γ_2 , and play b if the column player in Γ_2 ."

- $(S, \text{hare}, \text{rest})$ is an ESS if $\frac{R}{1 - \delta} > \frac{R}{1 - \delta \Delta}$ and $\frac{R}{1 - \delta} > \frac{r}{1 - \Delta}$. The former is always satisfied, but the latter depends on the exact values of R, r, δ & Δ . So, this strategy is an ESS whenever $\frac{R}{1 - \delta} > \frac{r}{1 - \Delta}$.
- $(SR, \text{hare}, \text{rest})$ is an ESS if $\frac{R}{1 - \delta \Delta} > \frac{R + \delta r}{1 - \delta \Delta}$ and $\frac{r}{1 - \Delta} < \frac{R}{1 - \delta \Delta}$. The former is never satisfied, so this strategy is never an ESS.
- $(H, \text{hare}, \text{rest})$ is an ESS if $\frac{r}{1 - \Delta} > 0$, which is always satisfied. So, this strategy is always an ESS.

21 Formulate IPD with TFT and sPav as a multi-stage game.

TFT	C D D	C D D	...	TFT	C C ...	sPav	D C C C ...
sPav	D D C	D D C	...	TFT	C C ...	sPav	D C C C ...

$S < P < R < T$

Γ_1	C	D
(TFT)	R + $\delta\Gamma_1$	S + $\delta\Gamma_2^{(col)}$
(sPav)	T + $\delta\Gamma_2^{(row)}$	P + $\delta\Gamma_4$

Γ_2	D (TFT)
D (sPav)	P + $\delta\Gamma_3^{(row)}$, P + $\delta\Gamma_3^{(col)}$

Γ_3	D (TFT)
C (sPav)	S + $\delta\Gamma_1$, T + $\delta\Gamma_1$

Γ_4	C (sPav)
C (sPav)	R + $\delta\Gamma_4$

(symmetric)

Strategy set (same for row & col players)

$$\{ (C, \binom{D}{D}), (C, \binom{C}{C}), (D, \binom{D}{D}), (D, \binom{C}{C}) \}$$

Note: $\binom{a}{b}$ is a conditional strategy wherein the player plays a if in the position of the row player and plays b if in the position of the column player.