

Evolution and the Theory of Games

Fall 2014

Exercises 5-8

5 Find all Nash equilibria (mixed and pure) of the Hawk-Dove game for $R > C$ and for $R < C$.

	H	D
H	$\frac{R-C}{2}, \frac{R-C}{2}$	$R, 0$
D	$0, R$	$\frac{R}{2}, \frac{R}{2}$

Pure strategy Nash equilibria:

$R > C$

$$\pi_1(D, H) = 0 < \frac{R-C}{2} = \pi_1(H, H)$$

$$\pi_1(D, D) = \frac{R}{2} < R = \pi_1(H, D)$$

$$\pi_2(H, D) = 0 < \frac{R-C}{2} = \pi_2(H, H)$$

$$\pi_2(D, D) = \frac{R}{2} < R = \pi_2(D, H)$$

$\Rightarrow (H, H)$ is the only Nash equilibrium when $R > C$.

$R < C$

$$\pi_1(H, H) = \frac{R-C}{2} < 0 = \pi_1(D, H)$$

$$\pi_1(D, D) = \frac{R}{2} < R = \pi_1(H, D)$$

$$\pi_2(H, H) = \frac{R-C}{2} < 0 = \pi_2(H, D)$$

$$\pi_2(D, D) = \frac{R}{2} < R = \pi_2(D, H)$$

$\Rightarrow (D, H)$ + (H, D) are Nash equilibria when $R < C$.

Mixed strategy Nash equilibria:

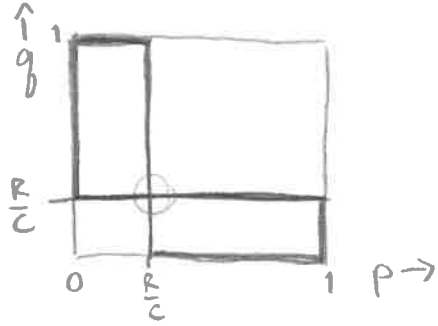
Let $x = (p, 1-p)$ + $y = (q, 1-q)$

$$\begin{aligned} \pi_1(x, y) &= pq\left(\frac{R-C}{2}\right) + p(1-q)R + (1-p)q \cdot 0 + (1-p)(1-q)\frac{R}{2} \\ &= pq\frac{R}{2} - pq\frac{C}{2} + pR - pqR + \frac{R}{2} - p\frac{R}{2} - q\frac{R}{2} + pq\frac{R}{2} \\ &= pq(R - \frac{C}{2} - R) + p\frac{R}{2} - q\frac{R}{2} + \frac{R}{2} \\ &= (1-q)\frac{R}{2} + p\frac{R}{2} - pq\frac{C}{2} \\ &= (1-q)\frac{R}{2} + \frac{1}{2}(R - qc)p \end{aligned}$$

If $R - qc > 0$ (i.e. $q < \frac{R}{c}$), x should play $p = 1$.

If $R - qc < 0$ (i.e. $q > \frac{R}{c}$), x should play $p = 0$.

If $R - qc = 0$ (i.e. $q = \frac{R}{c}$), x 's strategy makes no difference.



$$\begin{aligned} \pi_2(x, y) &= pg\left(\frac{R-C}{2}\right) + p(1-g) \cdot 0 + (1-p)gR + (1-p)(1-g)\frac{R}{2} \\ &= pg \cdot \frac{R}{2} - pg \cdot \frac{C}{2} + gR - pgR + \frac{R}{2} - p\frac{R}{2} - g\frac{R}{2} + pg\frac{R}{2} \\ &= \cancel{pgR} - \cancel{pgR} + g\frac{R}{2} - pg\frac{C}{2} + \frac{R}{2} - p\frac{R}{2} \\ &= (1-p)\frac{R}{2} + \left(\frac{R}{2} - p\frac{C}{2}\right)g \\ &= (1-p)\frac{R}{2} + \frac{1}{2}(R - pc)g \end{aligned}$$

- If $p > \frac{R}{C}$, y should play $g = 0$.
- If $p < \frac{R}{C}$, y should play $g = 1$.
- If $p = \frac{R}{C}$, y can choose any strategy.

So, we find the Nash equilibrium $\begin{cases} \hat{x} = (\frac{R}{C}, 1 - \frac{R}{C}) \\ \hat{y} = (\frac{R}{C}, 1 - \frac{R}{C}) \end{cases}$

Note that this equilibrium exists only when $\frac{R}{C} < 1$, i.e. when $\boxed{R < C}$.

Note: The top-left and bottom-right corners of the swastika diagram show the Nash equilibria $\begin{cases} \hat{x} = (0, 1) \\ \hat{y} = (1, 0) \end{cases}$ and $\begin{cases} \hat{x} = (1, 0) \\ \hat{y} = (0, 1) \end{cases}$, which correspond to the pure-strategy Nash equilibria, (D, H) and (H, D), respectively.

6 Suppose that (\hat{x}, \hat{y}) is a Nash equilibrium. Show that $\pi_i(x, \hat{y}) = \pi_i(\hat{x}, \hat{y})$ for every pure strategy x in the support of \hat{x} .

• Let $\hat{x} = (p_1, \dots, p_n)$, where p_i is the probability of playing pure strategy x_i .

• As (\hat{x}, \hat{y}) is a Nash equilibrium,

$$\pi_i(x_i, \hat{y}) \leq \pi_i(\hat{x}, \hat{y}) \quad \text{for } x_i \in X, i = 1, \dots, n$$

(Assume X has n elements.)

• Assume the inequality is strict, i.e. assume

$$\pi_i(x_i, \hat{y}) < \pi_i(\hat{x}, \hat{y}) \quad \text{for some } x_i \text{ in the support of } \hat{x},$$

and assume that $\pi_i(x_j, \hat{y}) = \pi_i(\hat{x}, \hat{y}) \quad \forall j \neq i$ (where x_j is in the support of \hat{x}).

$$\text{Then, } \pi_i(\hat{x}, \hat{y}) = \sum_j p_j \pi_i(x_j, \hat{y}) = p_i \pi_i(x_i, \hat{y}) + \sum_{j \neq i} p_j \pi_i(x_j, \hat{y})$$

$$< p_i \pi_i(\hat{x}, \hat{y}) + \sum_{j \neq i} p_j \pi_i(x_j, \hat{y}) = \pi_i(\hat{x}, \hat{y}),$$

which is impossible. Thus, for no x_i in the support of \hat{x} is it possible that $\pi_i(x_i, \hat{y}) < \pi_i(\hat{x}, \hat{y})$. So, $\pi_i(x, \hat{y}) = \pi_i(\hat{x}, \hat{y}) \quad \forall x$ in the support of \hat{x} , QED.

7 Show that every dominating strategy sol'n is a Nash equilibrium, but that the reverse is not necessarily true.

- $x \in X$ is dominated by $x' \in X$ if $\pi_i(x, y) \leq \pi_i(x', y) \quad \forall y \in Y$ with a strict inequality for at least one $y \in Y$.

- $(\hat{x}, \hat{y}) \in X \times Y$ is NE provided that

$$\begin{cases} \pi_1(x, \hat{y}) \leq \pi_1(\hat{x}, \hat{y}) \quad \forall x \in X \\ \pi_2(\hat{x}, y) \leq \pi_2(\hat{x}, \hat{y}) \quad \forall y \in Y \end{cases}$$

Assume that \exists DSS (x_j, y_k) which is not NE. This means that $\exists x_i \ni \pi_1(x_i, y_k) > \pi_1(x_j, y_k)$ [or that $\exists y_i \ni \pi_2(x_j, y_i) > \pi_2(x_j, y_k)$]. Without loss of generality, assume the former. In this case, row i of the payoff matrix can never be dominated by row j in the iterated removal process, meaning that this process cannot end with only the strategy pair (x_j, y_k) left. As this contradicts the assumption that (x_j, y_k) is a DSS, it is impossible that (x_j, y_k) be a DSS but not a NE.

For an example of a NE that is not a DSS, consider the H-D game with R,C. From exercise 5, we know that (H,D) + (D,H) are NE, but neither is a DSS.

	H	D
H	$\frac{1}{2}(R,C), \frac{1}{2}(R,C)$	R, 0
D	0, R	$\frac{R}{2}, \frac{R}{2}$

8 Show that if $x \in X$ is a strictly dominated pure strategy and $(\hat{x}, \hat{y}) \in X \times Y$ is NE, then x cannot be in the support of \hat{x} .

Show that this conclusion need not be true if x is only weakly dominated, using the following example payoff matrix.

	y_1	y_2	y_3
x_1	3,2	3,0	2,2
x_2	1,0	3,3	0,3
x_3	0,2	0,0	3,2

Suppose that x is in the support of \hat{x} . If x is strictly dominated by \bar{x} , then $\pi_i(x, y) < \pi_i(\bar{x}, y) \quad \forall y \in Y$. It follows immediately that $\pi_i(x, \hat{y}) < \pi_i(\bar{x}, \hat{y})$. Since x is in the support of \hat{x} , by the Bishop-Cannings theorem $\pi_i(x, \hat{y}) = \pi_i(\hat{x}, \hat{y}) < \pi_i(\bar{x}, \hat{y})$, which contradicts that (\hat{x}, \hat{y}) is NE. Thus, a strictly dominated pure strategy x cannot be in the support of \hat{x} .

In the above matrix, there are 3 NE:

(x_1, y_1) , (x_2, y_2) + (x_3, y_3) .

Notice that y_2 is weakly dominated by y_3 despite being in the support of a NE.

(y_1 is also weakly dominated by y_3 , and x_2 is weakly dominated by x_1)